



4. Semantic Processing and Attributed Grammars

Semantic Processing



The parser checks only the *syntactic* correctness of a program

Tasks of semantic processing

- **Checking context conditions**
 - Declaration rules
 - Type checking
- **Symbol table handling**
 - Maintaining information about declared names
 - Maintaining information about types
 - Maintaining scopes
- **Invocation of code generation routines**

Semantic actions are integrated into the parser.
We describe them with *attributed grammars*

Semantic Actions

So far, we have just analyzed the input

Number = digit {digit}.

the parser checks if the input is syntactically correct (in this example *Number* is not viewed as part of the lexical structure of the language)

Now, we also translate it (semantic processing)

e.g.: we want to count the digits in the number

Number =	
digit	(. int n = 1; .)
{ digit	(. n++; .)
}	
.	(. System.out.println(n); .)

syntax

semantics

semantic actions

- arbitrary Java statements between (. and .)
- are executed by the parser at the position where they occur in the grammar

"translation" here:

123 ⇒ 3

4711 ⇒ 4

9 ⇒ 1

Attributes

Syntax symbols can return values (sort of output parameters)

`digit <↑val>` *digit* returns its numeric value (0..9) as an output attribute

Attributes are useful in the translation process

e.g.: we want to compute the value of a number

```
Number      (. int val, n; .)
= digit <↑val>
  { digit <↑n>   (. val = 10 * val + n; .)
  }
              (. System.out.println(val); .)
.
```

"translation" here:

"123" ⇒ 123

"4711" ⇒ 4711

"9" ⇒ 9



Input Attributes

Nonterminal symbols can have also input attributes

(parameters that are passed from the "calling" production)

Number $\langle \downarrow \text{base}, \uparrow \text{val} \rangle$

base: number base (e.g. 10 or 16)

val: returned value of the number

Example

```
Number  $\langle \downarrow \text{base}, \uparrow \text{val} \rangle$   (. int base, val, n; .)
= digit  $\langle \uparrow \text{val} \rangle$ 
  { digit  $\langle \uparrow \text{n} \rangle$                 (. val = base * val + n; .)
  }.
```



Attributed Grammars

Notation for describing translation processes

consist of three parts

1. Productions in EBNF

```
IdentList = ident {"," ident}.
```

2. Attributes (parameters of syntax symbols)

```
ident<↑name>
```

```
IdentList<↓type>
```

output attributes (*synthesized*):

input attributes (*inherited*):

yield the translation result

provide context from the caller

3. Semantic actions

```
(. ... arbitrary Java statements ... .)
```

Example

ATG for processing declarations

```

VarDecl                (. Struct type; .)
= Type <↑type>
  IdentList <↓type>
  ";" .
  
```

```

IdentList <↓type>      (. Struct type; String name; .)
= ident <↑name>        (. Tab.insert(name, type); .)
  { "," ident <↑name>  (. Tab.insert(name, type); .)
  } .
  
```

This is translated to parsing methods as follows

```

private static void VarDecl() {
  Struct type;
  type = Type();
  IdentList(type);
  check(semicolon);
}
  
```

ATGs are shorter and more readable than parsing methods

```

private static void IdentList(Struct type) {
  String name;
  check(ident); name = t.val;
  Tab.insert(name, type);
  while (sym == comma) {
    scan();
    check(ident); name = t.val;
    Tab.insert(name, type);
  }
}
  
```

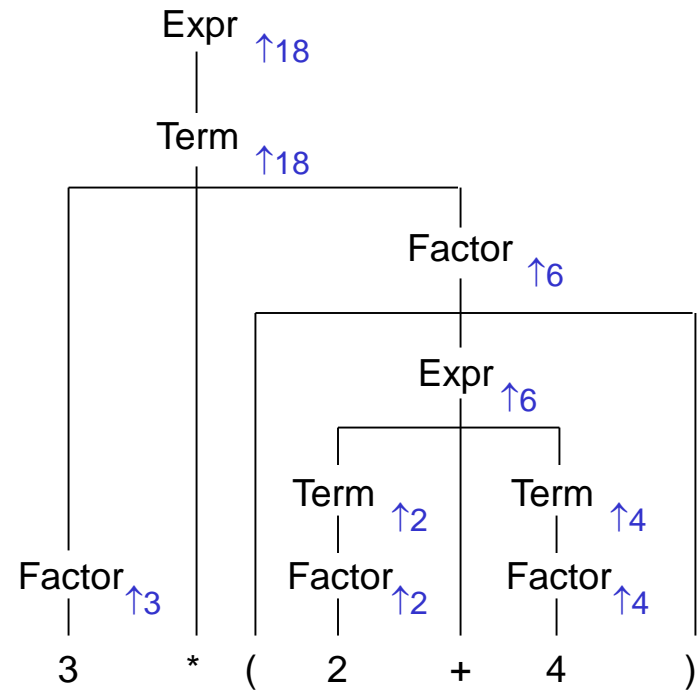
Example: Processing of Constant Expressions

input: 3 * (2 + 4)
 desired result: 18

```
Expr <↑val>      (. int val, val1; .)
= Term <↑val>
  { "+" Term <↑val1>  (. val = val + val1; .)
  | "-" Term <↑val1>  (. val = val - val1; .)
  }.
```

```
Term <↑val>      (. int val, val1; .)
= Factor <↑val>
  { "*" Factor <↑val1> (. val = val * val1; .)
  | "/" Factor <↑val1> (. val = val / val1; .)
  }
```

```
Factor <↑val>    (. int val, val1; .)
= number        (. val = t.numVal; .)
| "(" Expr <↑val> ")"
```



Transforming an ATG into a Parser

Production

```

Expr <↑val>          (. int val, val1; .)
= Term <↑val>
  { "+" Term <↑val1> (. val = val + val1; .)
  | "-" Term <↑val1> (. val = val - val1; .)
  }.

```

Parsing method

```

private static int Expr() {
    int val, val1;
    val = Term();
    for (;;) {
        if (sym == plus) {
            scan();
            val1 = Term();
            val = val + val1;
        } else if (sym == minus) {
            scan();
            val1 = Term();
            val = val - val1;
        } else break;
    }
    return val;
}

```

input attributes ⇒ parameters
output attribute ⇒ function value
(if there are multiple output attributes
encapsulate them in an object)

semantic actions ⇒ embedded Java code

Terminal symbols have no input attributes.

In our form of ATGs they also have no output attributes,
but their value can be obtained from *t.val* or *t.numVal*.



Example: Sales Statistics

ATGs can also be used in areas other than compiler construction

Example: given a file with sales numbers

```
File    = {Article}.
Article = Code {Amount} ";".
Code    = number.
Amount = number.
```

Whenever the input is syntactically structured
ATGs are a good notation to describe its processing

Input for example:

```
3451  2 5 3 7 ;
3452  4 8 1 ;
3453  1 1 ;
...
```

Desired output:

```
3451  17
3452  13
3453   2
...
```

ATG for the Sales Statistics



```
File                                (. int code, amount; .)
= { Article <↑code, ↑amount>          (. print(code + " " + amount); .)
  }.

Article <↑code, ↑amount>             (. int code, x, amount = 0; .)
= Number <↑code>
  { Number <↑x>                        (. amount += x; .)
  }
  ";\n";

Number <↑x>                          (. int x; .)
= number                               (. x = t.numVal; .)
.
```

Parser code

```
private static void File() {
    while (sym == number) {
        ArtInfo a = Article();
        print(a.code + " " + a.amount);
    }
}
```

```
class ArtInfo {
    int code, amount;
}
```

```
private static ArtInfo Article() {
    ArtInfo a = new ArtInfo();
    a.amount = 0;
    a.code = Number();
    while (sym == number) {
        int x = Number();
        a.amount += x;
    }
    check(semicolon); return a;
}
```

```
private static int Number() {
    check(number);
    return t.numVal;
}
```

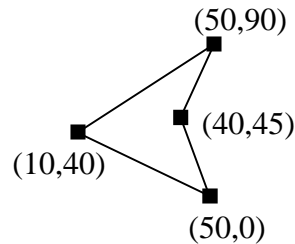
terminal symbols

number
semicolon
eof

Example: Image Description Language



described by:



```
POLY  
  (10,40)  
  (50,90)  
  (40,45)  
  (50,0)  
END
```

input syntax:

```
Polygon = "POLY" Point {Point} "END".  
Point = "(" number "," number ")".
```

We want a program that reads the input and draws the polygon

```
Polygon      (. Pt p, q; .)  
= "POLY"  
  Point<↑p>    (. Turtle.start(p); .)  
  { Point<↑q>  (. Turtle.move(q); .)  
  }  
  "END"        (. Turtle.move(p); .)  
.  
  
Point<↑p>    (. Pt p; int x, y; .)  
= "(" number  (. x = t.numVal; .)  
  "," number  (. y = t.numVal; .)  
  ")"         (. p = new Pt(x, y); .)  
.
```

We use "Turtle Graphics" for drawing

Turtle.start(p); sets the turtle (pen) to point *p*
Turtle.move(q); moves the turtle to *q*
drawing a line

```
class Pt {  
  int x;  
  int y;  
  Pt (int x, int y) { this.x = x; this.y = y; }  
}
```

Example: Transform Infix to Postfix Expressions



Arithmetic expressions in infix notation are to be transformed to postfix notation

$$3 + 4 * 2 \Rightarrow 3 4 2 * +$$

$$(3 + 4) * 2 \Rightarrow 3 4 + 2 *$$

```

Expr
= Term
  { "+" Term (. print("+"); .)
  | "-" Term (. print("-"); .)
  }.

Term
= Factor
  { "*" Factor (. print("*"); .)
  | "/" Factor (. print("/"); .)
  }.

Factor
= number (. print(t.numVal); .)
| "(" Expr ")".
    
```

